Dense graphs and a conjecture for tree-depth

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Joint work with John Sinkovic (University of Waterloo)

Tree-depth

(aka (vertex) ranking number, ordered coloring number, ...)



Tree-depth td(G): the smallest number of labels needed in a labeling of the vertices of *G* such that every path with equally-labeled endpoints also has a higher label. (Here, td(G) = 4)

Equivalently, the minimum number of vertex deletion steps needed to delete all of *G*, where in each step at most one vertex is deleted from each connected component.



















Theorem

If G contains H as a minor, then $td(G) \ge td(H)$.



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NOTE: A substructure of G forces G to have a higher tree-depth than it might otherwise have.

Call a graph **critical** if every proper minor has a smaller tree-depth. (k-critical = critical, with tree-depth k)

Critical graphs for small tree-depths (Dvořák–Giannopoulou–Thilikos, '09, '12)



• What structural properties must critical graphs possess?

Structural properties of k-critical graphs



Conjectures

If G is k-critical, then

- $|V(G)| \le 2^{k-1}$ [Dvořák–Giannopoulou–Thilikos, '09, '12]
- $\Delta(G) \leq k-1$

[B-Sinkovic, '16]

An approach to the max degree conjecture

Any vertex with the smallest label has neighbors with distinct labels.



If in some optimal labeling of *G* a vertex of maximum degree receives the label 1, then $\Delta(G) \leq td(G) - 1$.





Conjectures If *G* is *k*-critical, then

• *G* has an optimal labeling where some vertex with maximum degree is labeled 1.



Conjectures If *G* is *k*-critical, then

• For **each** vertex *v* of *G*, there is an optimal labeling where *v* receives label 1.



Conjectures If *G* is *k*-critical, then

• For each vertex v of G, there is an optimal labeling where v is the **unique** vertex receiving label 1.

1-uniqueness

Given a graph G, we say that a vertex v is **1-unique** if there is an optimal labeling of G where v is the only vertex receiving label 1.



We say a graph G is 1-unique if each vertex of G is 1-unique.

Observation

If G is 1-unique, then $\Delta(G) \leq td(G) - 1$.

Conjecture

If *G* is critical, then *G* is 1-unique.

M. D. Barrus (URI)

Plausibility: 1-uniqueness as a type of criticality?

Theorem (B-Sinkovic, 2016)

If G is a 1-unique graph, then

- deleting any vertex of G lowers the tree-depth, and
- contracting any edge of G lowers the tree-depth.

Hence G has a connected spanning subgraph that is td(G)-critical.

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NOTE: We get a td(G)-critical graph by deleting <u>some</u> (possibly none) of the edges of G.

Plausibility: small tree-depths



Conjecture

• If *G* is *k*-critical, then *G* is 1-unique.

True for $k \leq 4...$

Another application of 1-uniqueness



Theorem (B–Sinkovic, 2016)

Graphs G constructed with k-critical "appendages" L_i and an ℓ -critical "core" H...

- ...have tree-depth $k + \ell 1$ if H and all the L_i are critical;
- ...are critical if H is also 1-unique;
- ...are 1-unique if H and all the L_i are 1-unique;

Another application of 1-uniqueness



Theorem (B–Sinkovic, 2016)

Graphs G constructed with k-critical "appendages" L_i and an ℓ -critical "core" H...

…have order at most 2^{td(G)-1} if |V(H)| ≤ 2^{ℓ-1} and |V(L_i)| ≤ 2^{k-1} for all i.

Some notes



Observations

• Max degree conjecture always true, graph order bound <u>often</u> true (at least!) if all critical graphs are 1-unique.

Some notes



Observations

- Max degree conjecture <u>always</u> true, graph order bound <u>often</u> true (at least!) if all critical graphs are 1-unique.
- Most of the critical graphs seem to be sparse...

Of course, *K_n* is always dense, *n*-critical...

Curiosity: For which (critical?) graphs G is tree-depth close to n(G)?

High tree-depth

Let *G* be a graph with n(G) = n.

• td(G) = n iff G is $\{2K_1\}$ -free.



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• $td(G) \ge n-1$ iff G is $\{3K_1, 2K_2\}$ -free.

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Label the $3K_1$ the same; distinct labels on all other vertices.

$$1 ext{ td}(G) \le 1 + (n-3) = n-2.$$

EXAMPLE 2 Label the $2K_2$ with two labels; distinct labels on all other vertices.

$$td(G) \le 2 + (n-4) = n-2.$$

2

1

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 $td(G) \ge n-1$ iff G is $\{3K_1, 2K_2\}$ -free.

Low tree-depth forces induced subgraphs

If $td(G) \le n-2$ then fix an optimal labeling. Two vertices share labels with other vertices.

Case: Only one label appears on multiple vertices.

Case: Two labels appears on multiple vertices.



High tree-depth

Let *G* be a graph with n(G) = n.

• td(G) = n iff G is $\{2K_1\}$ -free.



- $td(G) \ge n-1$ iff G is $\{3K_1, 2K_2\}$ -free.
- $td(G) \ge n-2$ iff G is $\{4K_1, 2K_2 + K_1, P_3 + K_2, 2K_3\}$ -free.
- Conjectured: $td(G) \ge n-3$ iff *G* is { $5K_1, 2K_2 + 2K_1, P_3 + K_2 + K_1, 2K_3 + K_1, 2P_3, S_4 + K_2, 3K_2, diamond + K_3, C_4 + K_3, paw + K_3, P_4 + K_3, 2K_4$ }-free.

High tree-depth is hereditary

Theorem

Given a nonnegative integer k, let \mathcal{F}_k denote the set of minimal elements, under the induced subgraph ordering, of all graphs H for which $n(H) - td(H) \ge k + 1$. For all graphs G, the graph G satisfies $td(G) \ge n(G) - k$ if and only if G is \mathcal{F}_k -free.

Back to criticality and 1-uniqueness

Theorem

If G is an n-vertex critical graph and $td(G) \ge n - 1$, then G is 1-unique.

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Lemma

 $td(G) \ge n-1$ iff G is $\{3K_1, 2K_2\}$ -free.



Theorem (B–Sinkovic, 2016)

A graph G is 1-unique if and only if every star-clique transformation lowers the tree-depth.

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Outline: Easy to verify if td(G) = n, so assume td(G) = n - 1. Let *H* be a graph obtained from *G* via a star-clique transformation at *v*.

- *H* is still $\{3K_1, 2K_2\}$ -free.
- td(H) < td(G) unless *H* is a complete graph.
- G vu has tree-depth n 2; subgraphs induced force H to not be complete.

Interesting families of dense(ish) graphs

The Andrasfai graphs

And(k) has vertex set $\{0, \ldots, 3k - 2\}$.

Edges join vertices whose difference is 1 modulo 3. (The graph is a circulant graph.)

- And(k) is triangle-free.
- Every maximal independent set is a vertex neighborhood.
- And(k) is k-connected, k-regular, and has independence number k.



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- And(k) is triangle-free.
- Every maximal independent set is a vertex neighborhood.
- And(*k*) is *k*-connected, *k*-regular, and has independence number *k*.
- td(And(k)) = 2k and td(And(k) v) = 2k 1
- And(k) and And(k) v are both 1-unique.
- And(k) and And(k) v are both tree-depth-critical.

Separating 1-uniqueness from criticality

Conjecture

If G is critical, then G is 1-unique.

Theorem (B–Sinkovic, 2016)

If G is a 1-unique graph, then

- deleting any vertex of G lowers the tree-depth, and
- contracting any edge of G lowers the tree-depth.

Hence G has a connected spanning subgraph that is td(G)-critical.

NOTE: We get a td(G)-critical graph by deleting <u>some</u> (possibly none) of the edges of G. (How many?)

Another dense family: cycle complements



For all $n \ge 5$, td $\left(\overline{C_n}\right) = n - 1$ and $\overline{C_n}$ is 1-unique.

For $n \ge 8$, $\overline{C_n}$ is not critical.

But we can delete edges from $\overline{C_n}$ to reach an (n-1)-critical spanning subgraph...how many edges can it take?

1-unique can be far from critical

For all
$$n \ge 5$$
, td $\left(\overline{C_n}\right) = n - 1$ and $\overline{C_n}$ is 1-unique.

For n = 4k and $k \ge 2$, let G_n be obtained from $\overline{C_n}$ by deleting every other antipodal chord.

For $n = 4k \ge 8$, G_n is a $\{3K_1, 2K_2\}$ -free spanning subgraph having k fewer edges than $\overline{C_n}$.

(Note that G_n need not be critical. Empirically, it seems that $\overline{C_n}$'s nearest spanning critical subgraph differs from it by an increasing number of edges.)



Resolution

Conjecture

For any *k*, if *G* is *k*-critical, then *G* is 1-unique.

Known true for

$$k = 1, 2, 3, 4,$$
 $n - 1, n$

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Conjecture

For any *k*, if *G* is *k*-critical, then *G* is 1-unique.

Known true for

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 $n - 1, n$

False for

$$k = 5, 6, \dots, n-2$$

Counterexample



A 7-vertex critical graph G with td(G) = 5 = n(G) - 2 that is not 1-unique!

Counterexamples



For $t \ge 2$, form H_t by subdividing all edges incident with a vertex of K_{t+2} .

Here,
$$n(H_t) = 2t + 3$$

 $td(H_t) = t + 3 = n - t$

 H_t is critical.

The vertex at the subdivided edges' center cannot receive the unique 1—a star-clique transformation at this vertex yields $K_{t+1} \Box K_2$, which has tree-depth $\lceil (3/2)(t+1) \rceil \ge t+3$.



Finding counterexamples



Computer search using SageMath's graph database and functions

- Iterate through proper colorings of a graph.
- Identify colorings with a unique lowest label; identify 1-unique (or nearly 1-unique) graphs.
- Use 1-uniqueness to produce candidates for criticality testing.

Ongoing questions



- Using non-1-unique critical graphs to construct larger critical graphs.
- Empirically, it appears each k-critical non-1-unique graph yields a (k 1)-critical graph when the sole problem vertex is removed. Is this always the case?
- How to prove/disprove that critical graphs satisfy Δ(G) ≤ td(G) − 1?

Thank you!