

The A_4 -structure of a graph

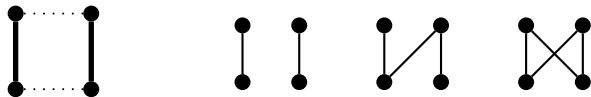
Michael D. Barrus

Department of Mathematics
Brigham Young University

Graphs and Matrices Seminar
September 24, 2012

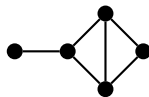
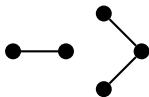
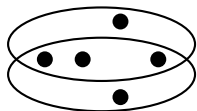
The A_4 -Structure

Alternating 4-cycle (A_4)

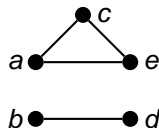
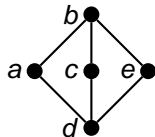
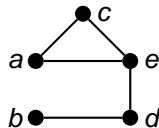
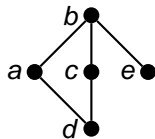
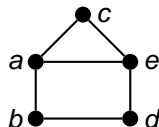
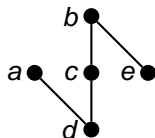
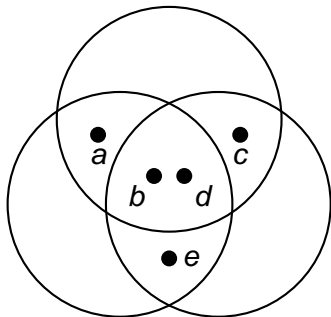


A_4 -structure H of a graph G

$$V(H) = V(G), \quad E(H) = \{A \subseteq V(G) : G[A] \cong 2K_2 \text{ or } C_4 \text{ or } P_4\}$$



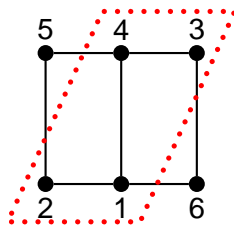
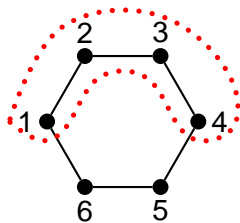
The A_4 -Structure



The P_4 -Structure of a Graph

Chvátal, 1984

$$V(H) = V(G), \quad E(H) = \{A \subseteq V(G) : G[A] \cong P_4\}$$



Theorem (Reed, 1987)

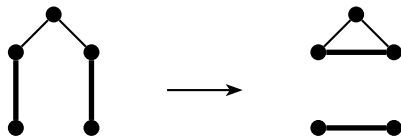
Let G and H be two graphs with isomorphic P_4 -structures. Then G is perfect if and only if H is perfect.

P_4 -Classes

Reprinted from A. Brandstädt and V. B. Le, Split-perfect graphs: characterization and algorithmic use, SIAM J. Discrete Math. 17(3), 341-360.

Motivation: Degree Sequences

2-switches



Theorem (Fulkerson–Hoffman–McAndrew, 1965)

$\deg(G) = \deg(H)$ iff 2-switches transform G into H .

Are there any A_4 -structure/degree sequence connections?

Motivation: Graph Classes

- Threshold graphs
 $\{2K_2, C_4, P_4\}$ -free
- Matrogenic graphs
Vertex sets of A_4 's are circuits of a matroid on V .
- Matroidal graphs
Edge sets of A_4 's are circuits of a matroid on E .

Motivation: Graph Classes

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No 5 vertices contain exactly 2 or 3 edges of the A_4 -structure.

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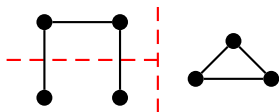
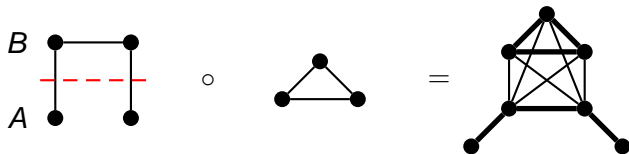
No 5 vertices contain more than 1 edge of the A_4 -structure.

Can the A_4 -structure be used to characterize other interesting classes?

A Graph Operation

Definition (Tyshkevich–Chernyak, 1978).

Given a split graph G with stable set A and clique B , and an arbitrary graph H , define the *composition* $(G, A, B) \circ H$ to be graph formed by adding to $G + H$ the edges in $\{uv : u \in B, v \in V(H)\}$.



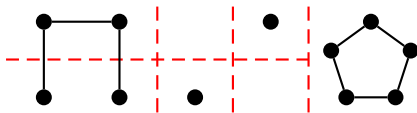
Canonical Decomposition

Theorem (Tyshkevich–Chernyak, 1978; Tyshkevich, 2000)

Every graph F can be represented as a composition

$$F = (G_k, A_k, B_k) \circ \cdots \circ (G_1, A_1, B_1) \circ F_0$$

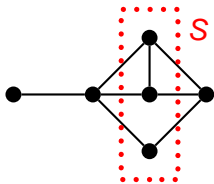
of indecomposable components. Here the (G_i, A_i, B_i) are indecomposable splitted graphs and F_0 is an indecomposable graph. This decomposition is unique up to isomorphism of components.



Modules and P_4 s

Definition.

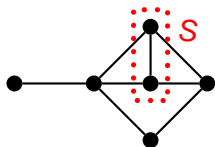
A *module* is a vertex subset S such that each vertex outside S either dominates S or is isolated from S .



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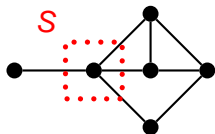
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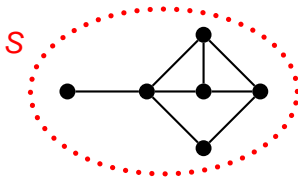
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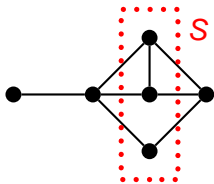
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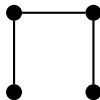
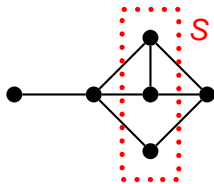
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Theorem

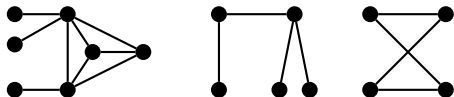
- An induced P_4 intersects a module in exactly 0, 1, or 4 vertices.
- (Seinsche, 1974) In a graph G every induced subgraph on at least 3 vertices contains a nontrivial module iff G is P_4 -free.

P_4 -Structures and Decomposition

Primeval Decomposition Theorem (Jamison–Olariu, 1995)

For any graph $G = (V, E)$ precisely one of the following conditions holds:

- (i) G is disconnected.
- (ii) \overline{G} is disconnected.
- (iii) The P_4 -structure of G is connected.
- (iv) There exists a P_4 -component hooked up to the rest of G in a special way.

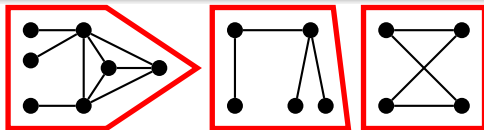


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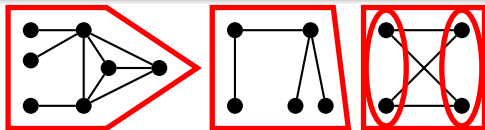


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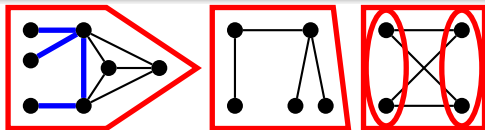


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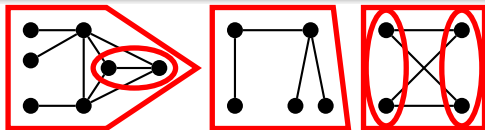


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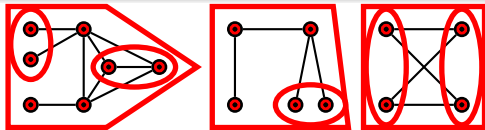


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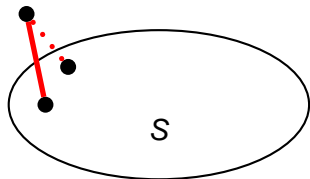


A_4 -Analogues

Definition.

A *module* S is a vertex subset such that no alternating path of length 2 begins and ends in S and has its midpoint outside S .

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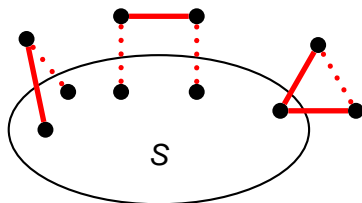


A_4 -Analogues

Definition.

Define a **strict module** to be a vertex subset S such that no (**possibly closed**) alternating path of length 2 or 3 begins and ends in S and has its midpoints outside S .

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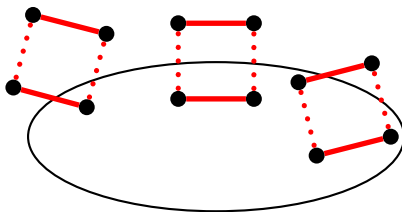
This is equivalent to not having alternating paths of *any* length begin and end in S .

A_4 -Analogues

Proposition

An A_4 intersects a strict module in exactly 0 or 4 vertices.

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Proposition

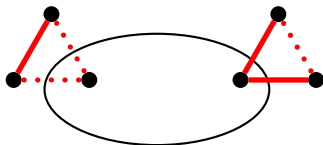
In a graph G every induced subgraph on at least 2 vertices has a nontrivial strict module if and only if G is A_4 -free, i.e., threshold.

Strict Modules and Graph Structure

Proposition

The vertices which dominate a strict module form a clique, and those which are nonadjacent to the strict module form an independent set.

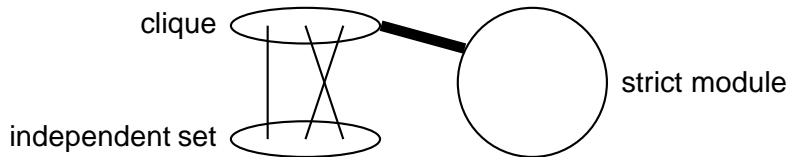
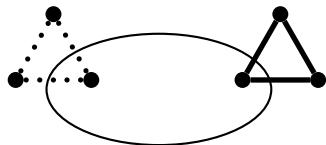
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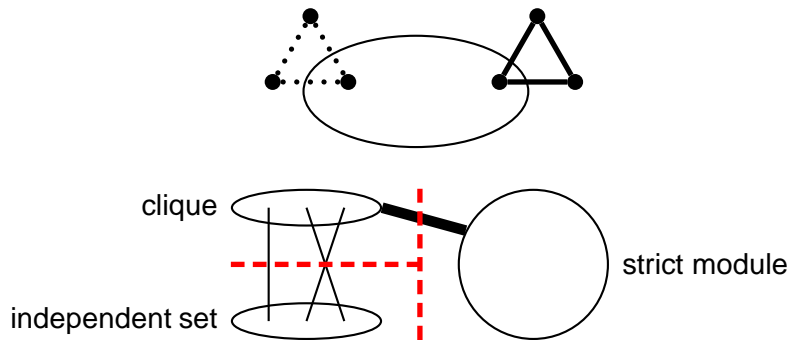
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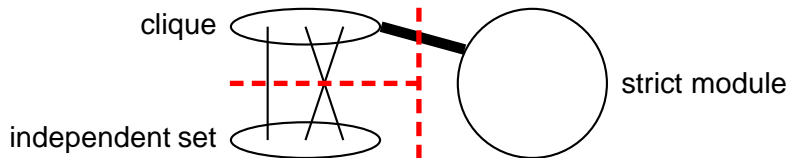
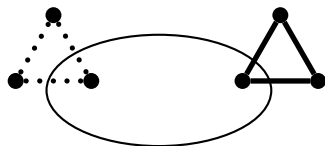
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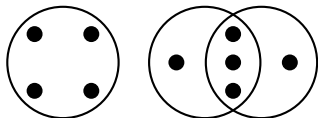
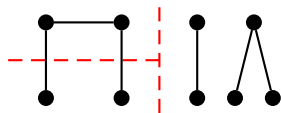
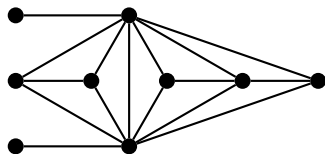


“Strict modular decomposition” = canonical decomposition

Indecomposable Graphs

Theorem

A graph is indecomposable in the canonical decomposition if and only if its A_4 -structure is connected.



► Proof later?

A Degree Sequence Connection

Theorem (Tyshkevich, 1980?, 2000)

An n -vertex graph with degree sequence d is decomposable if and only if there exists nonnegative integers p and q such that

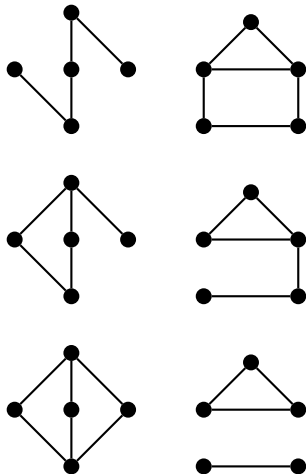
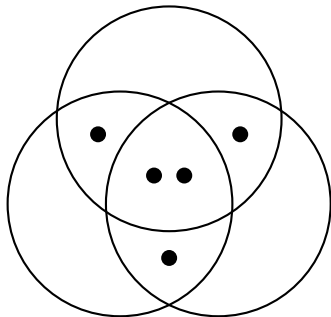
$$0 < p + q < n, \quad \sum_{i=1}^p d_i = p(n - q - 1) + \sum_{i=n-q+1}^n d_i.$$

Corollary

If two graphs have the same degree sequence, then their A_4 -structures have the same number and orders of components.

Obtaining All Realizations

Given an A_4 -structure, how do we generate all graphs realizing it?



Obtaining Other Realizations: Decomposable Graphs

Substitutions and transpositions preserve A_4 -structure.



The rightmost A_4 -component may only be transposed if it has a split realization.

Which graphs have the same A_4 -structure as a split graph?

A_4 -Separable Graphs

Observation

A graph G is A_4 -split, i.e., it has the same A_4 -structure as a split graph, iff each of its indecomposable component is A_4 -split.

A graph G is A_4 -separable if we can partition $V(G)$ into two sets so that each A_4 can be drawn with both edges and both nonedges spanning the divide.



A_4 -separable \implies A_4 -split

A_4 -Balanced Graphs with the BRP

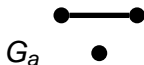
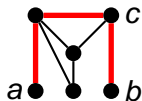
A graph G is A_4 -balanced if we can partition $V(G)$ into two sets so that each set contains two vertices of each A_4 . An A_4 -balanced graph has the bipartite restriction property if for each v , the graph G_v is bipartite.



A_4 -split \implies A_4 -balanced with the bipartite restriction property

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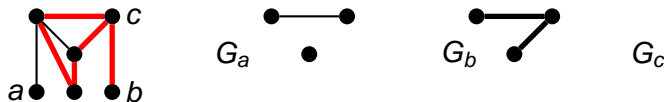
G_b

G_c

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A_4 -Balanced Graphs with the BRP

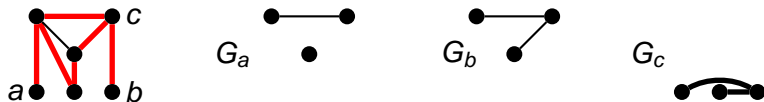
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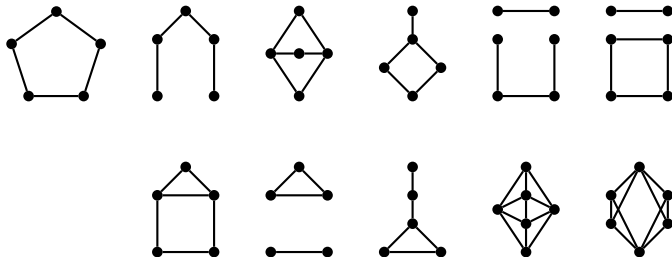
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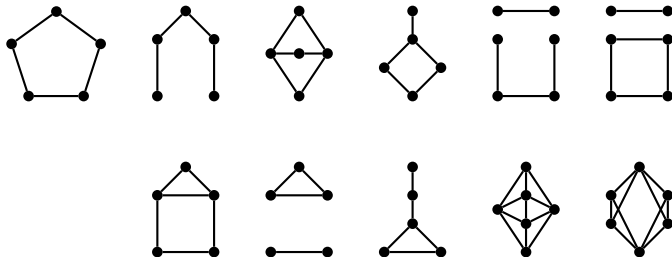
Forbidden subgraphs

The following graphs are not A_4 -balanced or do not have the BRP:



Forbidden subgraphs

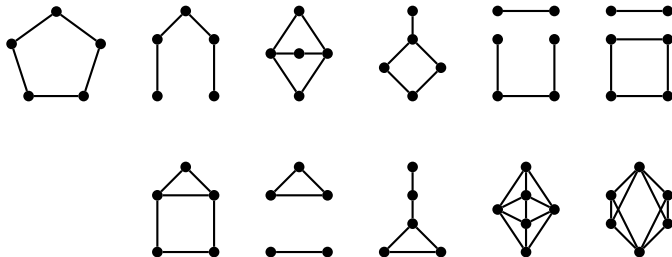
Say G induces none of the forbidden graphs:



G disconnected \implies Each component is $\{K_3, C_4, P_4\}$ -free

Forbidden subgraphs

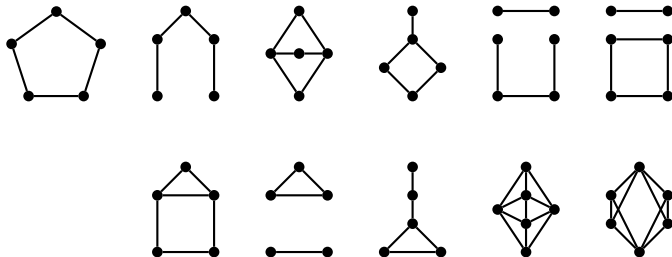
Say G induces none of the forbidden graphs:



G disconnected \implies Each component is a star

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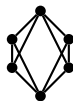
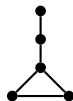
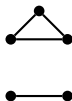
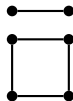
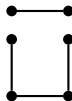
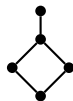
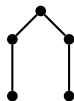
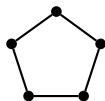


G connected, co-connected $\implies G$ is split.

Forbidden subgraphs

Say G induces none of the forbidden graphs:

\mathcal{F}



G A_4 -balanced, has BRP $\implies G$ is \mathcal{F} -free $\implies G$ is split, or G or \overline{G} is a forest of stars

Completing the chain



G split, or G or \overline{G} a forest of stars $\implies G$ A_4 -separable

A_4 -Split Graphs

Theorem

For an indecomposable graph G with A_4 -structure H , the following are equivalent:

- (i) G is A_4 -split.
- (ii) H is balanced and satisfies the bipartite restriction property.
- (iii) G is $\{C_5, P_5, \text{house}, K_2 + K_3, K_{2,3}, P, \overline{P}, K_2 + P_4, P_4 \vee 2K_1, K_2 + C_4, 2K_2 \vee 2K_1\}$ -free.
- (iv) G is split, or G or \overline{G} is a disjoint union of stars.
- (v) G is A_4 -separable.



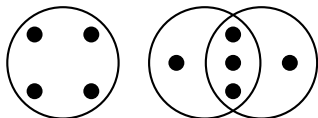
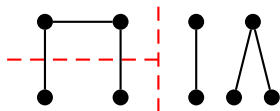
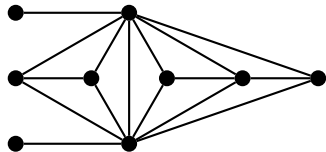
Left to Do

- Graph classes, especially A_4 - and P_4 -balanced graphs, and the A_4 -analogues of the (q, t) graphs (Threshold = $(4, 0)$, matroidal = $(5, 1)$, ...).
- Other A_4 -structure characteristics dependent only on degree sequence.
- Complete list of operations which suffice to link all realizations of an A_4 -structure.

Appendix

Theorem

A graph is indecomposable in the canonical decomposition if and only if its A_4 -structure is connected.

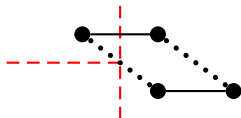


Beginnings

Lemma

The graphs $2K_2$, C_4 , and P_4 are all indecomposable. Therefore, connected A_4 -structure \implies indecomposable.

Forbidden:



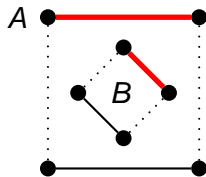
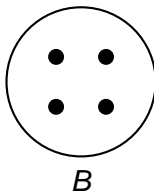
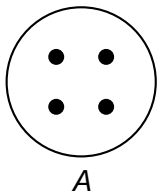
Lemma

In an indecomposable graph G with more than 1 vertex, every vertex belongs to an alternating 4-cycle.

Disjoint A_4 s

Lemma

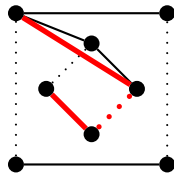
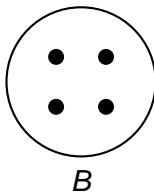
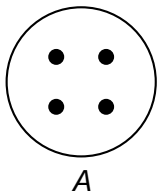
If A and B are disjoint alternating 4-cycles in G such that no third alternating cycle in G intersects each, then either A induces P_4 , with its interior vertices dominating B and the endpoints isolated from B (denote this by $A \rightarrow B$), or vice versa.



Disjoint A_4 s

Lemma

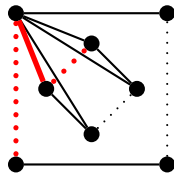
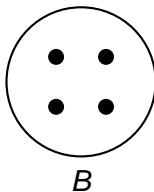
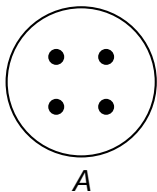
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Disjoint A_4 s

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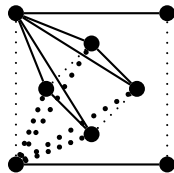
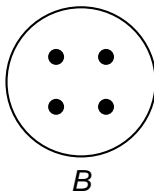
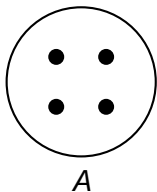
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Disjoint A_4 s

Lemma

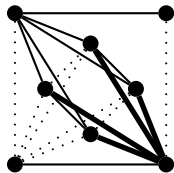
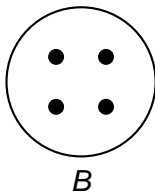
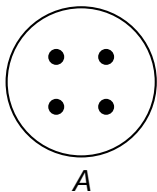
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Disjoint A_4 s

Lemma

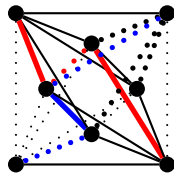
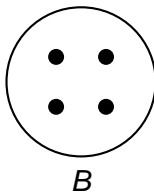
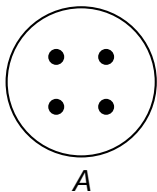
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Disjoint A_4 s

Lemma

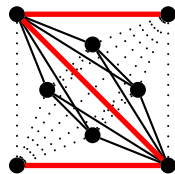
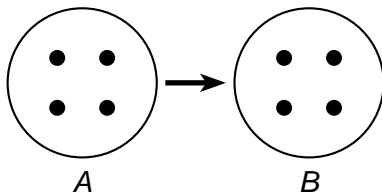
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Disjoint A_4 s

Lemma

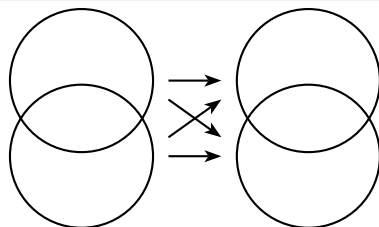
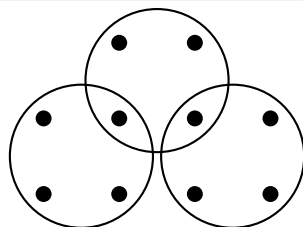
If A and B are disjoint alternating 4-cycles in G such that no third alternating cycle in G intersects each, then either A induces P_4 , with its interior vertices dominating B and the endpoints isolated from B (denote this by $A \rightarrow B$), or vice versa.



More on Disjoint A_4 s

Corollary

Any two vertices which both belong to induced $2K_2$'s or C_4 's have distance at most 3 in the A_4 -structure.



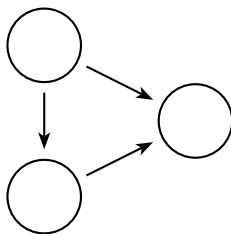
Lemma

The \rightarrow relation is consistent among pairs of A_4 s from different components of the A_4 -structure.

Putting It All Together

Lemma

The \rightarrow tournament on the A_4 -components of a graph is acyclic.



Having a source implies the graph is decomposable.

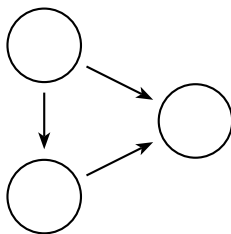
\therefore not A_4 -connected \implies decomposable.

[Return to talk](#)

Putting It All Together

Lemma

The \rightarrow tournament on the A_4 -components of a graph is acyclic.



Having a source implies the graph is decomposable.

$\therefore A_4$ -connected \iff indecomposable.