

Uniqueness in labelings of tree-depth-critical graphs

Michael D. Barrus

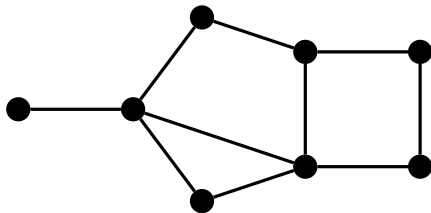
Department of Mathematics
University of Rhode Island

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Joint Mathematics Meetings 2017 • January 6, 2017

Joint work with John Sinkovic (University of Waterloo)

Tree-depth

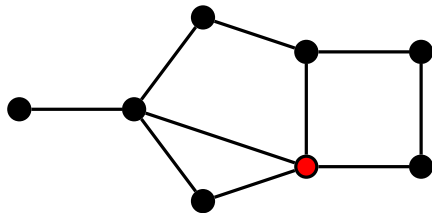
(aka (vertex) ranking number, ordered coloring number, ...)



Tree-depth $td(G)$: The minimum number of steps needed to delete all of G , where in each step at most one vertex is deleted from each connected component.

Tree-depth

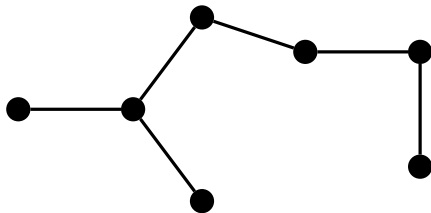
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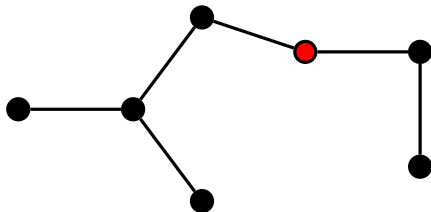
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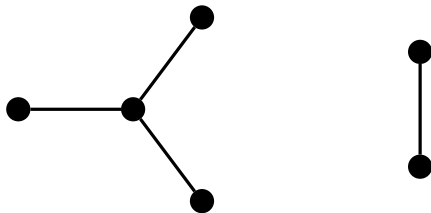
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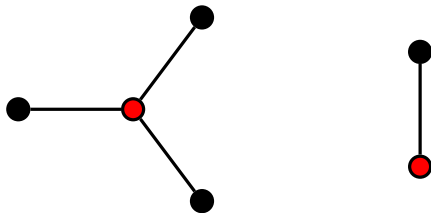
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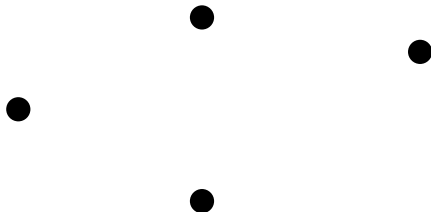
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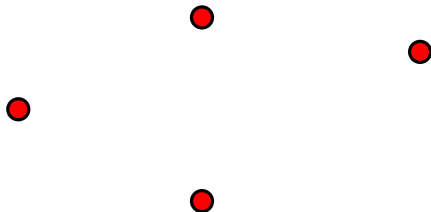
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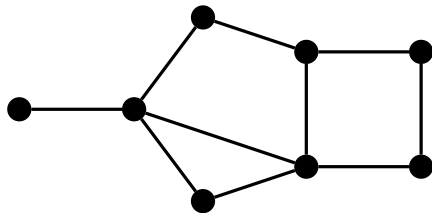
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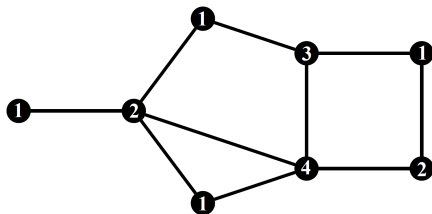
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Tree-depth $td(G)$: The minimum number of steps needed to delete all of G , where in each step at most one vertex is deleted from each connected component. (Here, $td(G) = 4$)

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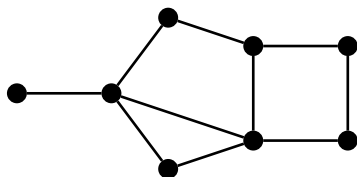
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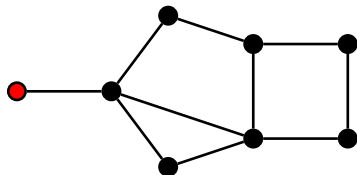
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Equivalently, the smallest number of labels needed in a labeling where every path with equal endpoints also has a higher label.

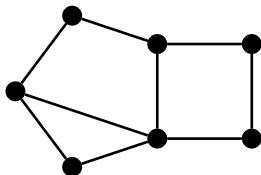
Tree-depth and criticality (with respect to minors)



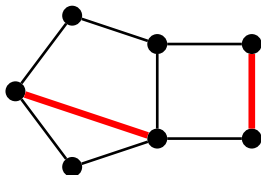
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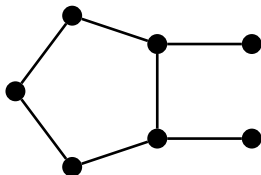
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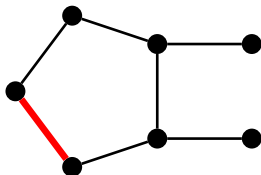
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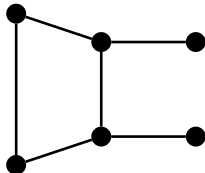
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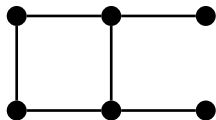
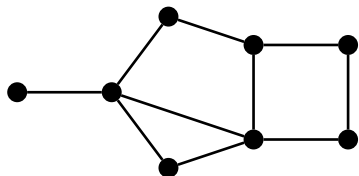
Tree-depth and criticality (with respect to minors)



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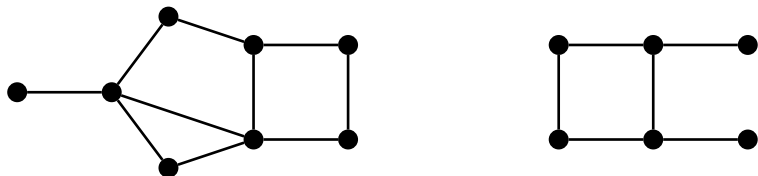
Tree-depth and criticality (with respect to minors)



Theorem

If G contains H as a minor, then $\text{td}(G) \geq \text{td}(H)$.

Tree-depth and criticality (with respect to minors)



Theorem

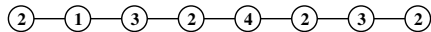
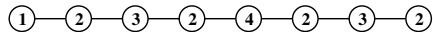
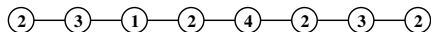
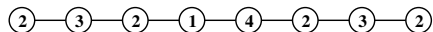
If G contains H as a minor, then $\text{td}(G) \geq \text{td}(H)$.

Idea: A graph is **critical** if every proper minor has lower tree-depth.

Critical minors determine the tree-depth.

New concept: 1-uniqueness

A graph is **1-unique** if for every vertex v there exists an optimal ranking where v is the only vertex assigned the label 1.



Why 1-uniqueness? It's plausible!

Critical graphs with small tree-depth: Dvořák–Giannopoulou–Thilikos, '09, '12

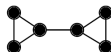
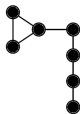
1:



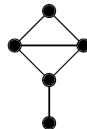
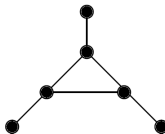
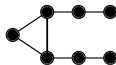
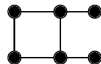
2:



3:



4:



5: 136 trees,
plus...

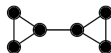
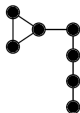
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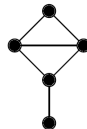
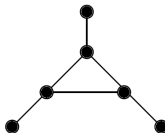
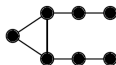
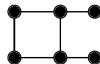
1: ●

2: ●—●

3: ●—●—●—●



4:



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(Ultimately hasty) conjecture: (B, Sinkovic, 2016)

All critical graphs are 1-unique!

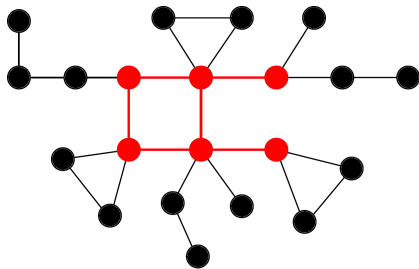
Why 1-uniqueness? It's plausible!

Just like critical graphs, 1-unique graphs are tree-depth-critical with respect to

- **vertex deletions,**
- **edge contractions,** and
- **deletions of edge cuts.**

Why 1-uniqueness? It's useful!

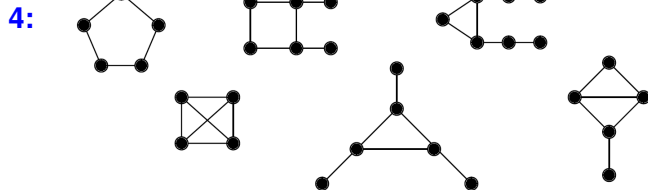
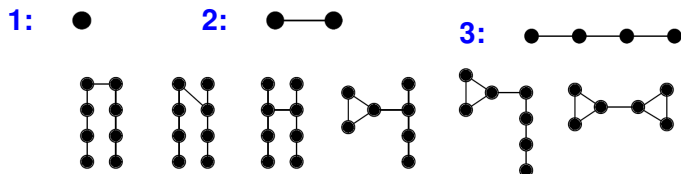
Hang k -minor-critical “appendages” off every vertex of an ℓ -minor-critical graph...



Theorem (B, Sinkovic, 2016)

Graphs constructed in this way are $(k + \ell - 1)$ -minor-critical if the appendages are 1-unique.

Why 1-uniqueness? It's useful!



5: 136 trees,
plus...

B, Sinkovic, 2016

Conjecture: If G is k -critical, then $\Delta(G) \leq k - 1$

Proposition: If G is 1-unique and $\text{td}(G) = k$, then $\Delta(G) \leq k - 1$.

Back to that conjecture

(B, Sinkovic, 2016, 2017+)

Is every critical graph 1-unique?

Yes for

- every k -critical graph for $k \in \{1, 2, 3, 4\}$,
- every critical tree,
- every critical cycle,
- every Andrasfai graph,
- ...

Back to that conjecture

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Theorem

If G is an n -vertex critical graph **and** $\text{td}(G) \geq n - 1$, then G is 1-unique.

Resolution

Conjecture

For any k , if G is k -critical, then G is 1-unique.

Known true for

$$k = 1, 2, 3, 4,$$

$$n - 1, n$$

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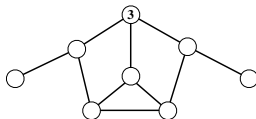
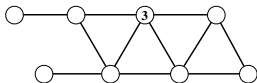
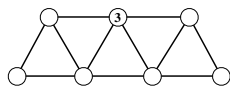
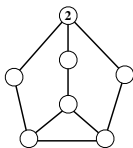
$$k = 1, 2, 3, 4,$$

$$n - 1, n$$

False for

$$k = 5, 6, \dots, n - 2$$

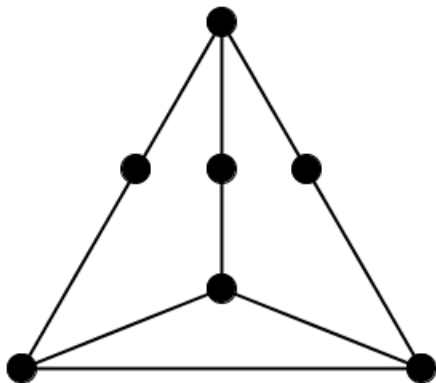
Finding counterexamples



Computer search using SageMath's graph database and functions

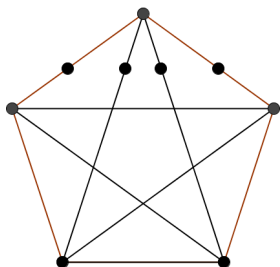
- Iterate through proper colorings of a graph.
- Identify colorings with a unique lowest label; identify 1-unique (or nearly 1-unique) graphs.
- Use 1-uniqueness to produce candidates for criticality testing.

Counterexample



A 7-vertex critical graph G with $\text{td}(G) = 5 = n(G) - 2$ that is not 1-unique!

Counterexamples



For $t \geq 2$, form H_t by subdividing all edges incident with a vertex of K_{t+2} .

Here, $\text{td}(H_t) = n(H_t) - t$.

The graph H_t is critical, but in **no** optimal ranking can the vertex at the subdivided edges' center receive the unique 1.

Conjecture

For any k , if G is k -critical, then G is 1-unique.

Known true for

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Resolution

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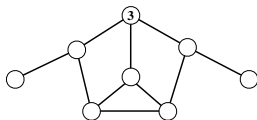
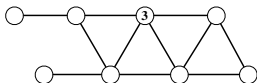
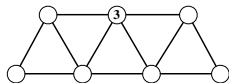
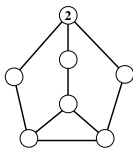
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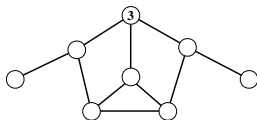
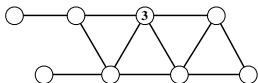
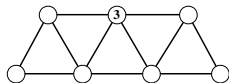
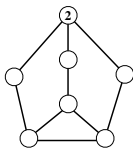
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Remaining questions



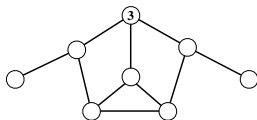
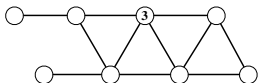
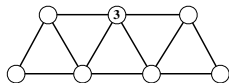
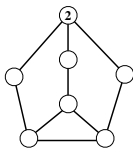
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Remaining questions



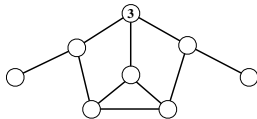
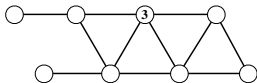
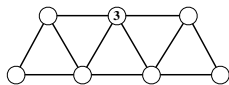
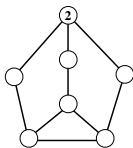
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- What fraction of critical graphs are 1-unique?

Remaining questions



- Empirically, it appears that when a k -critical non-1-unique graph has a unique “problem vertex,” deleting it yields a $(k - 1)$ -critical graph. Is this always the case?
- What fraction of critical graphs are 1-unique?
- How else, instead, can we prove/disprove that critical graphs satisfy
$$\Delta(G) \leq \text{td}(G) - 1?$$

Remaining questions



Thank you!