# Adjacency relationships forced by graph degree sequences

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## The question

$$d = (2, 2, 1, 1)$$

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### Which edges and non-edges are forced by the degree sequence?

M. D. Barrus (BYU)

**Threshold sequence:** a degree sequence having exactly one realization.



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All edges and non-edges are forced by the degree sequence.

## An answer

A clique is **demanding** if every vertex outside the clique has as many neighbors as possible in the clique.



Theorem

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## An answer

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Forced edges correspond exactly to three cases:







Forced adjacency relationships

**Threshold sequence:** a degree sequence having exactly one realization.



All edges and non-edges are forced by the degree sequence.

Erdős–Gallai inequalities for degree sequences  $(d_1, \ldots, d_n)$ :

$$\underbrace{\sum_{i \leq k} d_i}_{\text{LHS}(k)} \leq \underbrace{k(k-1) + \sum_{i > k} \min\{k, d_i\}}_{\text{RHS}(k)}$$

for all  $k \leq m = \max\{i : d_i \geq i - 1\}$ .

**Threshold graphs:** RHS(k) - LHS(k) = 0 for all  $k \in \{1, ..., m\}$ .

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#### Theorem

A sequence d has **no** forcible edges if and only if RHS(1) – LHS(1)  $\geq$  1 and RHS(k) – LHS(k)  $\geq$  2 for all  $k \in \{2, ..., m\}$ .

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Intersection envelope graph



Union envelope graph



### Observation

If G is threshold iff 
$$I_{d(G)} = G = U_{d(G)}$$
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#### Theorem

For any degree sequence d, both  $I_d$  and  $U_d$  are threshold graphs.



Partitions of 2m under the dominance order



Partitions of 2*m* under the dominance order

Threshold sequences: maximal degree sequences



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#### Theorem

If vertices i and j have a forcible adjacency relationship in realizations of d, then i and j have the same adjacency relationship for all degree sequences that majorize d.



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Threshold sequences: maximal degree sequences

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## Thank you!

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